Chapter 1: Introduction

<u>Our goal:</u>

- get "feel" and terminology
- more depth, detail later in course
- □ approach:
 - use Internet as example

Overview:

- what's the Internet
- what's a protocol?
- network edge
- network core
- access net, physical media
- Internet/ISP structure
- performance: loss, delay
- protocol layers, service models
- network modeling

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Chapter 1: roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models
- 1.8 History



Task Force

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What's the Internet: a service view

communication *infrastructure* enables distributed applications:

• Web, email, games, ecommerce, file sharing

communication services provided to apps:

- Connectionless unreliable
- connection-oriented reliable



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What's a protocol?

<u>human protocols:</u>

- "what's the time?"
- "I have a question"
- introductions

... specific msgs sent

... specific actions taken when msgs received, or other events

network protocols:

- machines rather than humans
- all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

What's a protocol?

a human protocol and a computer network protocol:



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A closer look at network structure:

network edge: applications and hosts

network core:

- o routers
- network of networks

access networks, physical media: communication links



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The network edge:

end systems (hosts):

- run application programs
- e.g. Web, email
- at "edge of network"

client/server model

 client host requests, receives service from always-on server • e.g. Web browser/server; email client/server

peer-peer model:

- minimal (or no) use of dedicated servers
- e.g. Gnutella, KaZaA

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Network edge: connection-oriented service

<u>Goal:</u> data transfer between end systems

- handshaking: setup (prepare for) data transfer ahead of time
 - Hello, hello back human protocol
 - set up "state" in two communicating hosts
- TCP Transmission Control Protocol
 - Internet's connectionoriented service

TCP service [RFC 793]

- reliable, in-order bytestream data transfer
 - loss: acknowledgements and retransmissions
- flow control:
 - sender won't overwhelm receiver
- congestion control:
 - senders "slow down sending rate" when network congested

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Network edge: connectionless service

<u>Goal</u>: data transfer between end systems o same as before!

- UDP User Datagram Protocol [RFC 768]:
 - connectionless
 - unreliable data transfer
 - o no flow control
 - no congestion control

App's using TCP:

 HTTP (Web), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

 streaming media, teleconferencing, DNS, Internet telephony

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The Network Core

- mesh of interconnected routers
- <u>the</u> fundamental question: how is data transferred through net?
 - circuit switching: dedicated circuit per call: telephone net
 - packet-switching: data sent thru net in discrete "chunks"



Network Core: Circuit Switching

End-end resources reserved for "call"

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required



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Network Core: Circuit Switching

network resources (e.g., bandwidth) divided into "pieces"

- pieces allocated to calls
- resource piece *idle* if not used by owning call (no sharing)
- dividing link bandwidth into "pieces"
 - o frequency division
 - \bigcirc time division



Network Core: Packet Switching

each end-end data stream divided into *packets*

- user A, B packets share network resources
- each packet uses full link bandwidth

resources used as needed



resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
 - Node receives complete packet before forwarding

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Packet switching versus circuit switching

Packet switching allows more users to use network!





Is packet switching a "slam dunk winner?"

- Great for bursty data
 - resource sharing
 - o simpler, no call setup
- Excessive congestion: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 6)

Packet-switching: store-and-forward



- Takes L/R seconds to transmit (push out) packet of L bits on to link with capacity of R bps
- Entire packet must arrive at router before it can be transmitted on next link: store and forward

delay = 3L/R

Example: L = 7.5 Mbits R = 1.5 Mbps delay = 15 sec

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Packet-switched networks: forwarding

- Goal: move packets through routers from source to destination
 - we'll study several path selection (i.e. routing) algorithms (chapter 4)

datagram network:

- destination address in packet determines next hop
- routes may change during session
- analogy: driving, asking directions

virtual circuit network:

- each packet carries tag (virtual circuit ID), tag determines next hop
- fixed path determined at *call setup time*, remains fixed thru call
- routers maintain per-call state

Network Taxonomy

