

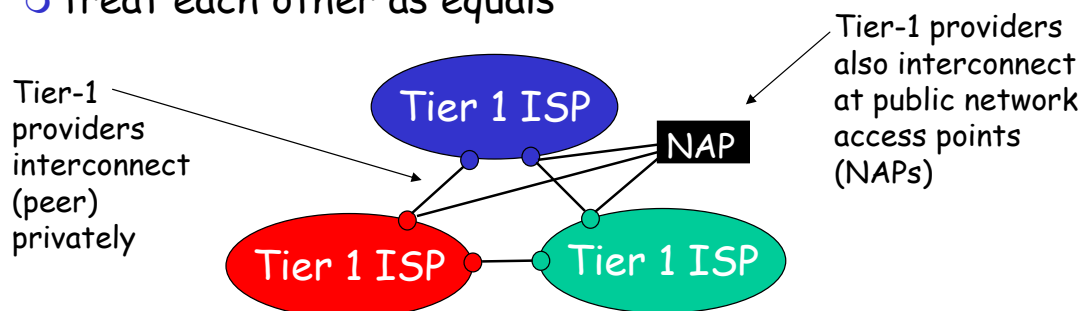
Chapter 1: roadmap

- 1.1 What *is* the Internet?
- 1.2 Network edge
- 1.3 Network core
- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
- 1.6 Delay & loss in packet-switched networks
- 1.7 Protocol layers, service models

Introduction 1-1

Internet structure: network of networks

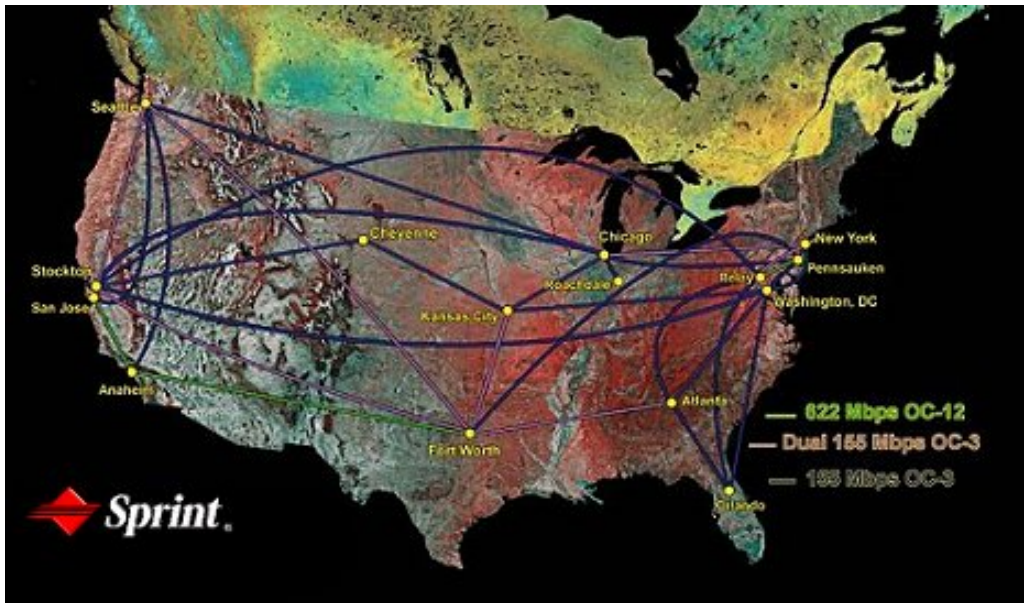
- roughly hierarchical
- at center: "tier-1" ISPs (e.g., UUNet, BBN/Genuity/level3, Sprint, AT&T, QWest), national/international coverage
 - treat each other as equals



Introduction 1-2

Tier-1 ISP: e.g., Sprint

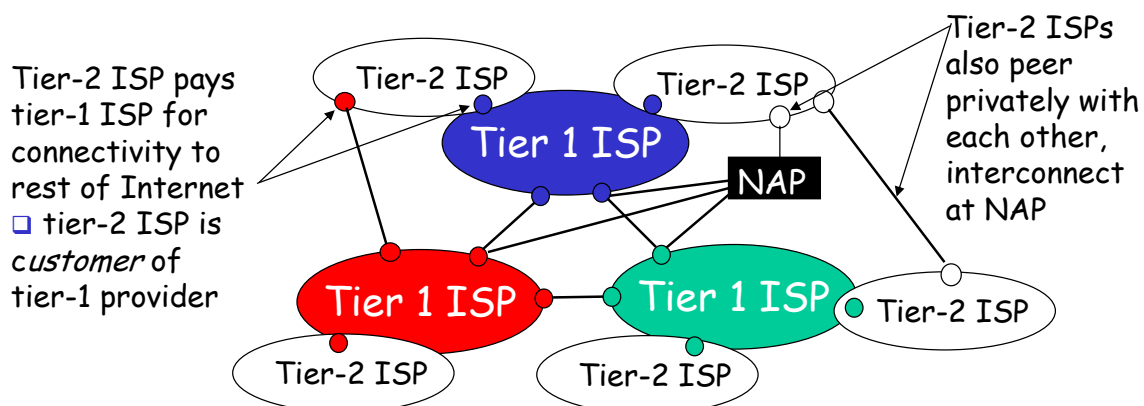
Sprint US backbone network



Introduction 1-3

Internet structure: network of networks

- "Tier-2" ISPs: smaller (often regional) ISPs
 - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

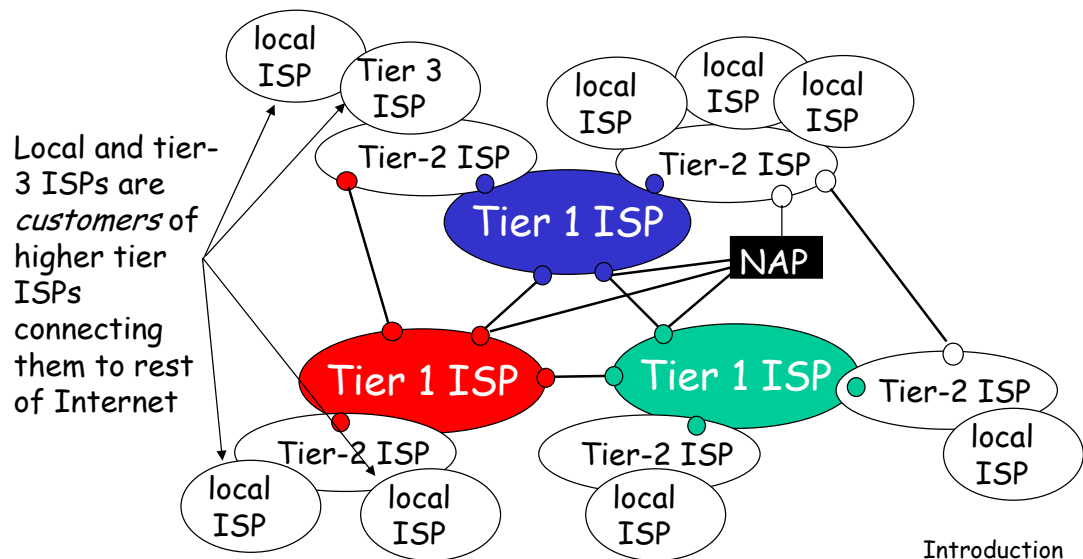


Introduction 1-4

Internet structure: network of networks

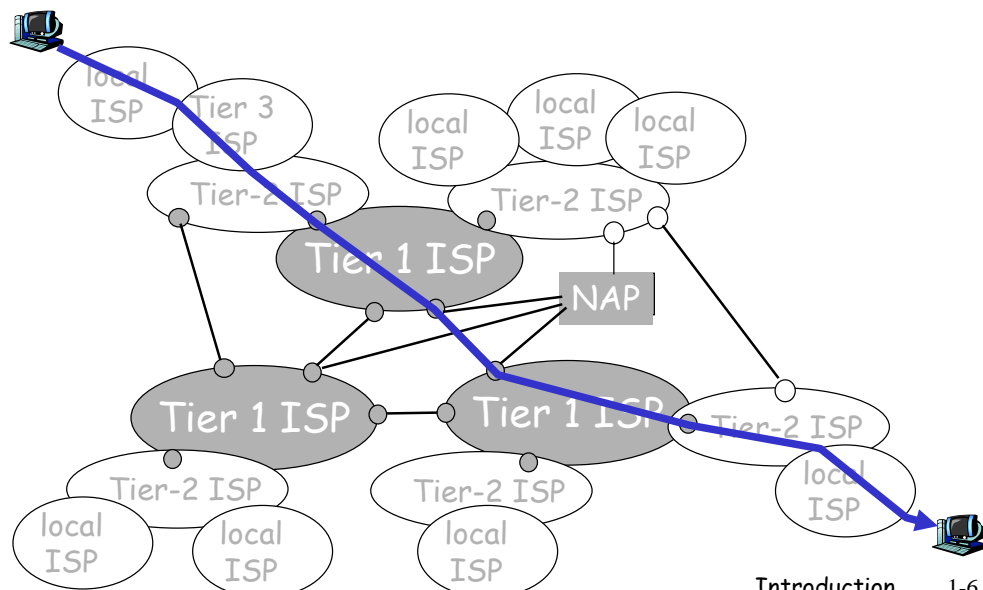
❑ "Tier-3" ISPs and local ISPs

- last hop ("access") network (closest to end systems)



Internet structure: network of networks

❑ a packet passes through many networks!



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1.6 Delay & loss in packet-switched networks

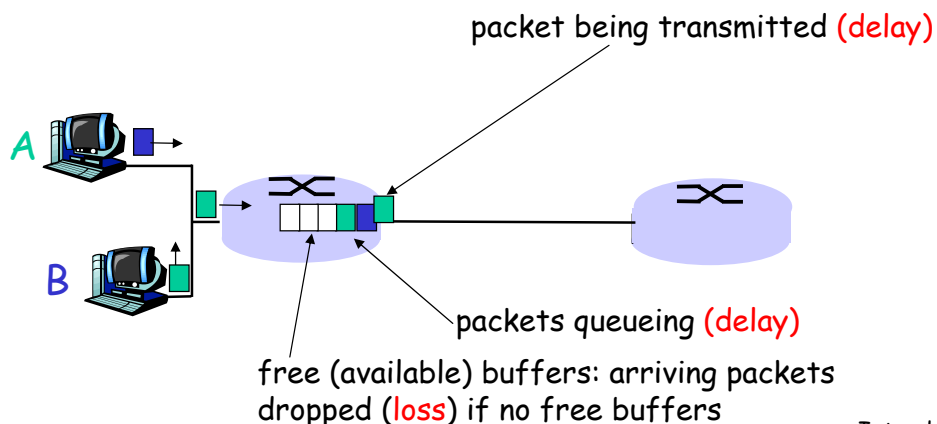
1.7 Protocol layers, service models

Introduction 1-7

How do loss and delay occur?

packets *queue* in router buffers

- ❑ packet arrival rate to link exceeds output link capacity
- ❑ packets queue, wait for turn



Introduction 1-8

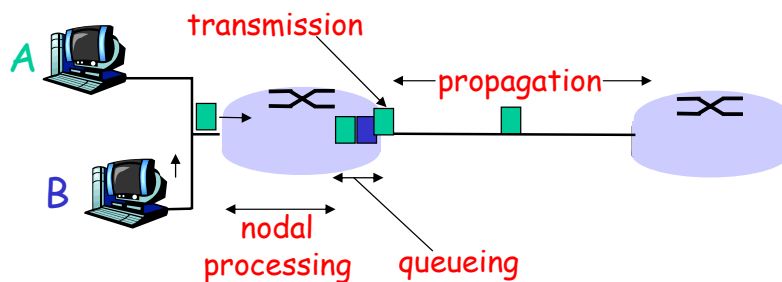
Four sources of packet delay

1. nodal processing:

- check bit errors
- determine output link

2. queueing

- time waiting at output link for transmission
- depends on congestion level of router



Delay in packet-switched networks

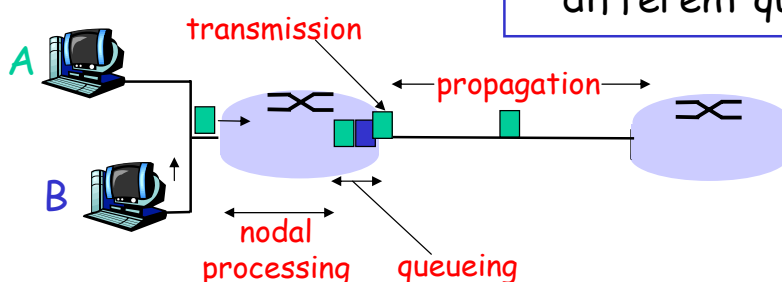
3. Transmission delay:

- R = link bandwidth (bps)
- L = packet length (bits)
- time to send bits into link = L/R

4. Propagation delay:

- d = length of physical link
- s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- propagation delay = d/s

Note: s and R are very different quantities!



Nodal delay

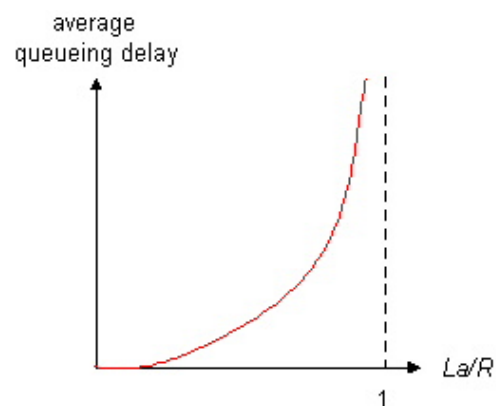
$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- d_{proc} = processing delay
 - typically a few microsecs or less
- d_{queue} = queuing delay
 - depends on congestion
- d_{trans} = transmission delay
 - $= L/R$, significant for low-speed links
- d_{prop} = propagation delay
 - a few microsecs to hundreds of msecs

Queueing delay (revisited)

- R = link bandwidth (bps)
- L = packet length (bits)
- a = average packet arrival rate

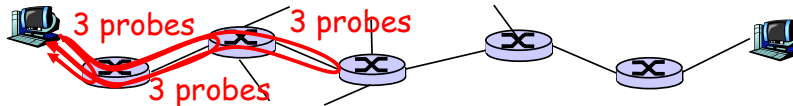
traffic intensity = La/R



- $La/R \sim 0$: average queueing delay small
- $La/R \rightarrow 1$: delays become large
- $La/R > 1$: more "work" arriving than can be serviced, average delay infinite!

"Real" Internet delays and routes

- ❑ What do "real" Internet delay & loss look like?
- ❑ Traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



Introduction 1-13

"Real" Internet delays and routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

Three delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu

```

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 ***
18 ***
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
  
```

trans-oceanic link

* means no response (probe lost, router not replying)

Introduction 1-14

Packet loss

- ❑ queue (aka buffer) preceding link in buffer has finite capacity
- ❑ when packet arrives to full queue, packet is dropped (aka lost)
- ❑ lost packet may be retransmitted by previous node, by source end system, or not retransmitted at all

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Protocol "Layers"

Networks are complex!

- ❑ many "pieces":
 - hosts
 - routers
 - links of various media
 - applications
 - protocols
 - hardware, software

Question:

Is there any hope of
organizing structure of
network?

Or at least our discussion
of networks?

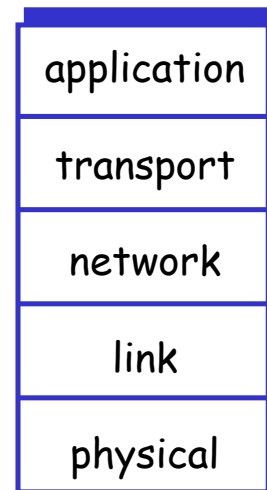
Why layering?

Dealing with complex systems:

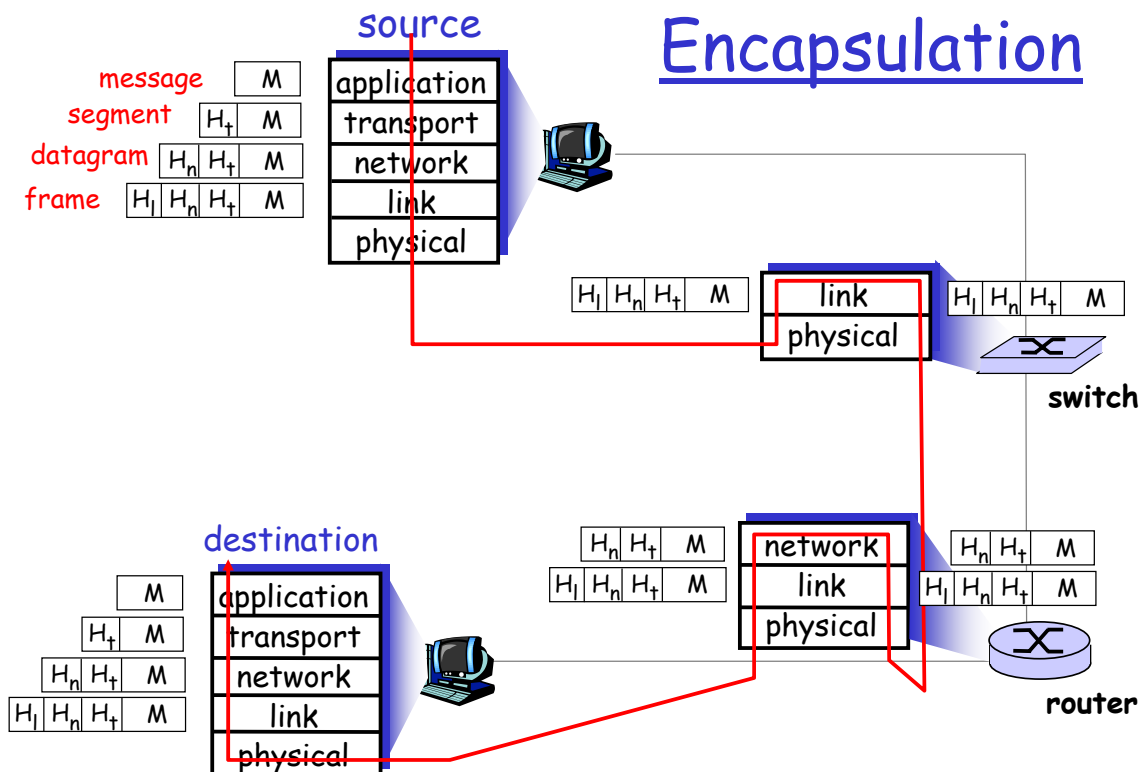
- ❑ explicit structure allows identification, relationship of complex system's pieces
 - layered **reference model** for discussion
- ❑ modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in one procedure doesn't affect rest of system
- ❑ layering considered harmful?

Internet protocol stack

- ❑ **application:** supporting network applications
 - FTP, SMTP, STTP
- ❑ **transport:** host-host data transfer
 - TCP, UDP
- ❑ **network:** routing of datagrams from source to destination
 - IP, routing protocols
- ❑ **link:** data transfer between neighboring network elements
 - PPP, Ethernet
- ❑ **physical:** bits "on the wire"



Introduction 1-19



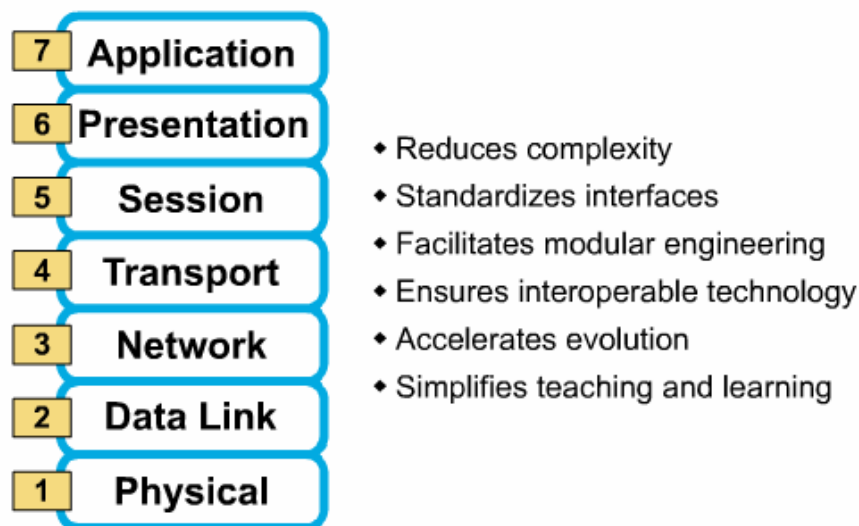
Introduction 1-20

OSI Reference Model

The International Standards Organization (ISO) proposal for the standardization of the various protocols used in computer networks (specifically those networks used to connect open systems) is called the Open Systems Interconnection Reference Model (1984), or simply the OSI model.

Introduction 1-21

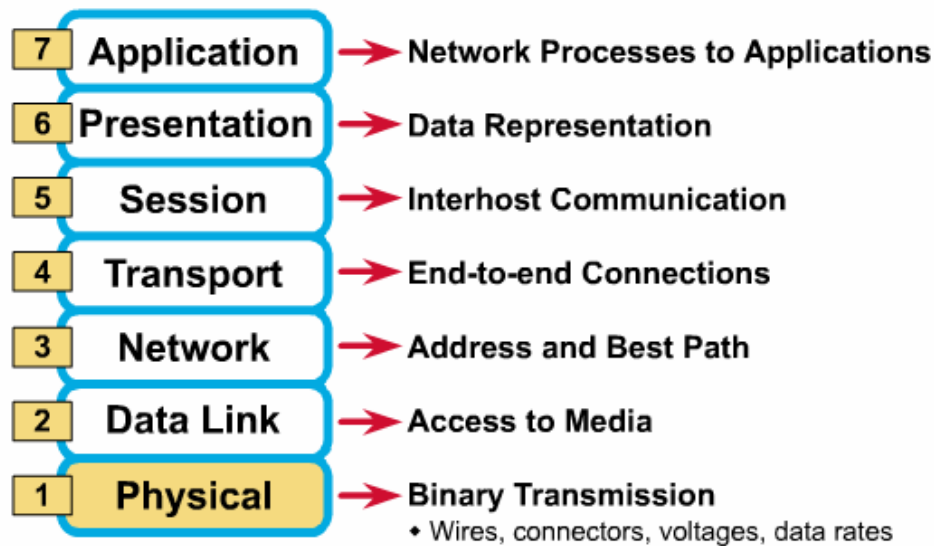
Why a Layered Model?



All People Seem To Need Data Processing

Introduction 1-22

Layers with Functions



Introduction 1-23

Introduction: Summary

Covered a "ton" of material!

- ❑ Internet overview
- ❑ what's a protocol?
- ❑ network edge, core, access network
 - packet-switching versus circuit-switching
- ❑ Internet/ISP structure
- ❑ performance: loss, delay
- ❑ layering and service models

You now have:

- ❑ context, overview, "feel" of networking
- ❑ more depth, detail *to follow!*

Introduction 1-24